# Optimizing Ruby's Memory Layout

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# How does Ruby manage memory?

## **RVALUE structure**

## RVALUE

**RClass** 

# **Ruby Object Structure**

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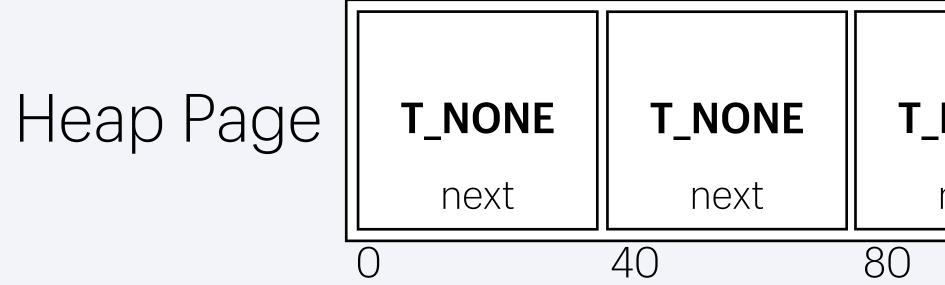
**RClass specific fields** 

## Heap page structure

- Heap pages are a container for a 16Kb memory region
- 409 slots per page

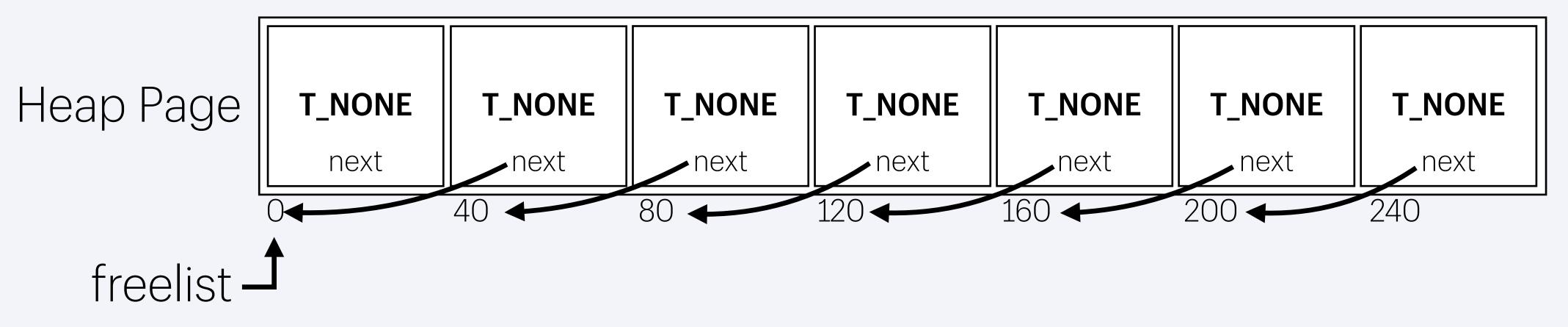
• All slots on the same page are contiguous. No gaps in addresses

## Heap page structure

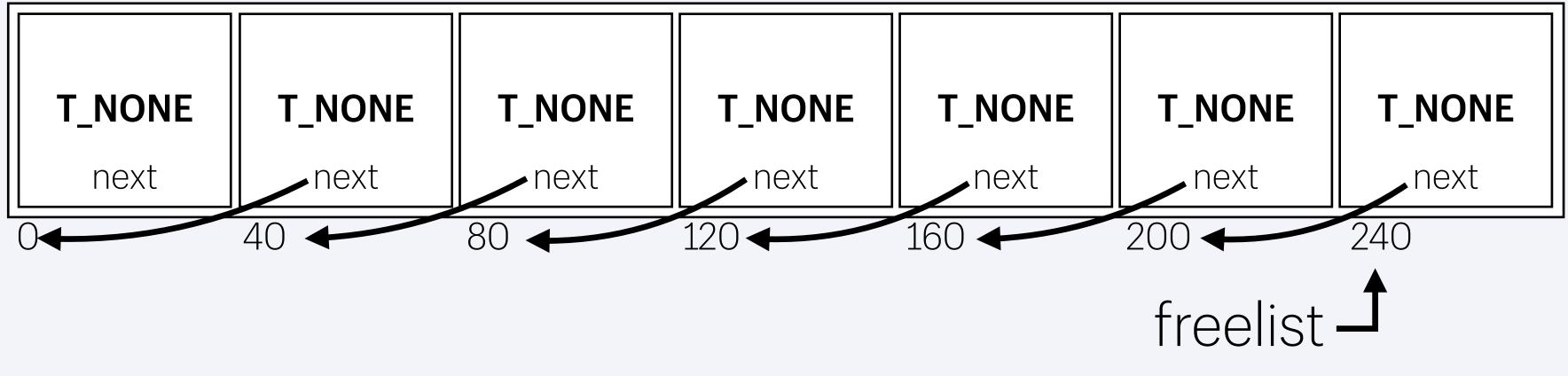


NONE	T_NONE	T_NONE	T_NONE	T_NONE
next	next	next	next	next
	120	160	200	240

# **Building the freelist**



## Allocating Ruby objects



# Heap Page

RClass	RHash	RString	RString	RArray	RString	RClass
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# How does Ruby's Garbage Collector work?

# Ruby's garbage collector

- Three phases:
  - Marking
  - Sweeping
  - Compaction (optional) •
- Stop-the-world garbage collection
- ٠

Execute Ruby code

Mark

Disclaimer: algorithms are simplified and some details are skipped Execute Ruby code Sweep Compact Time

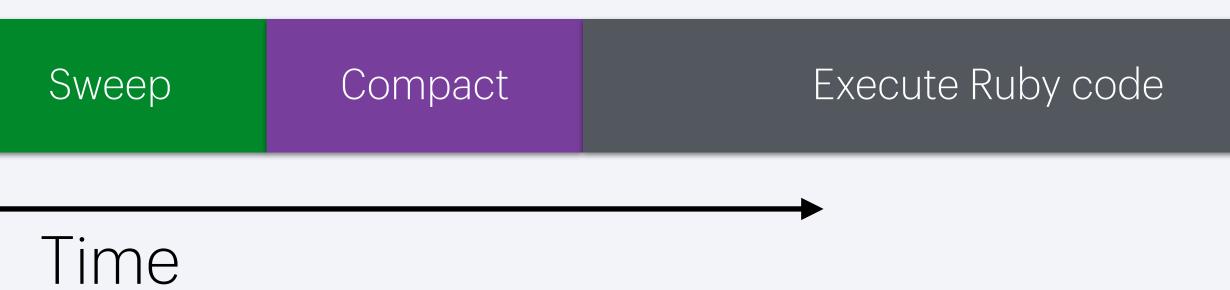
# Marking phase

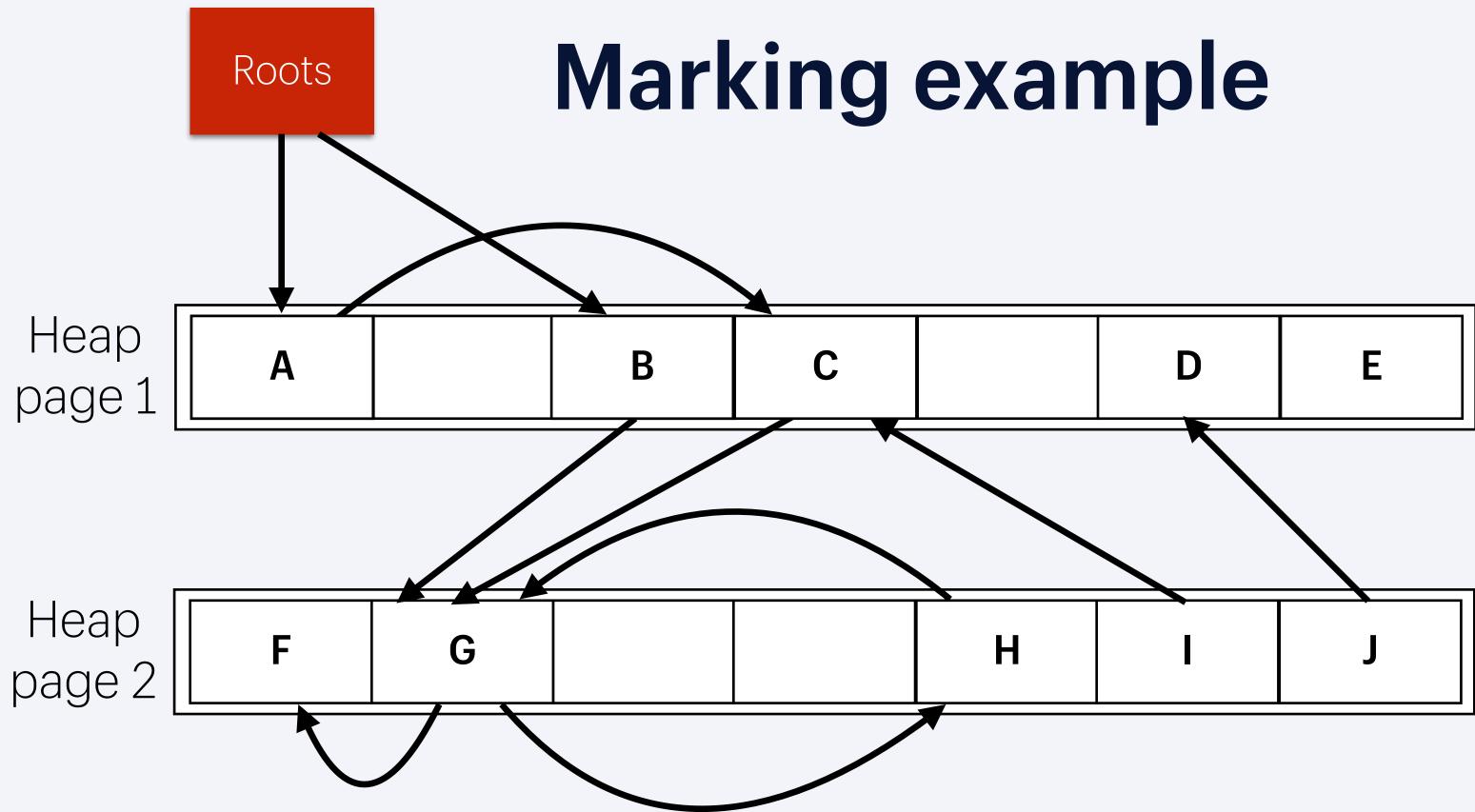
- Determines which Ruby objects are alive
- Push the object onto the mark stack when marked
- Recursively mark unmarked children of marked objects until empty



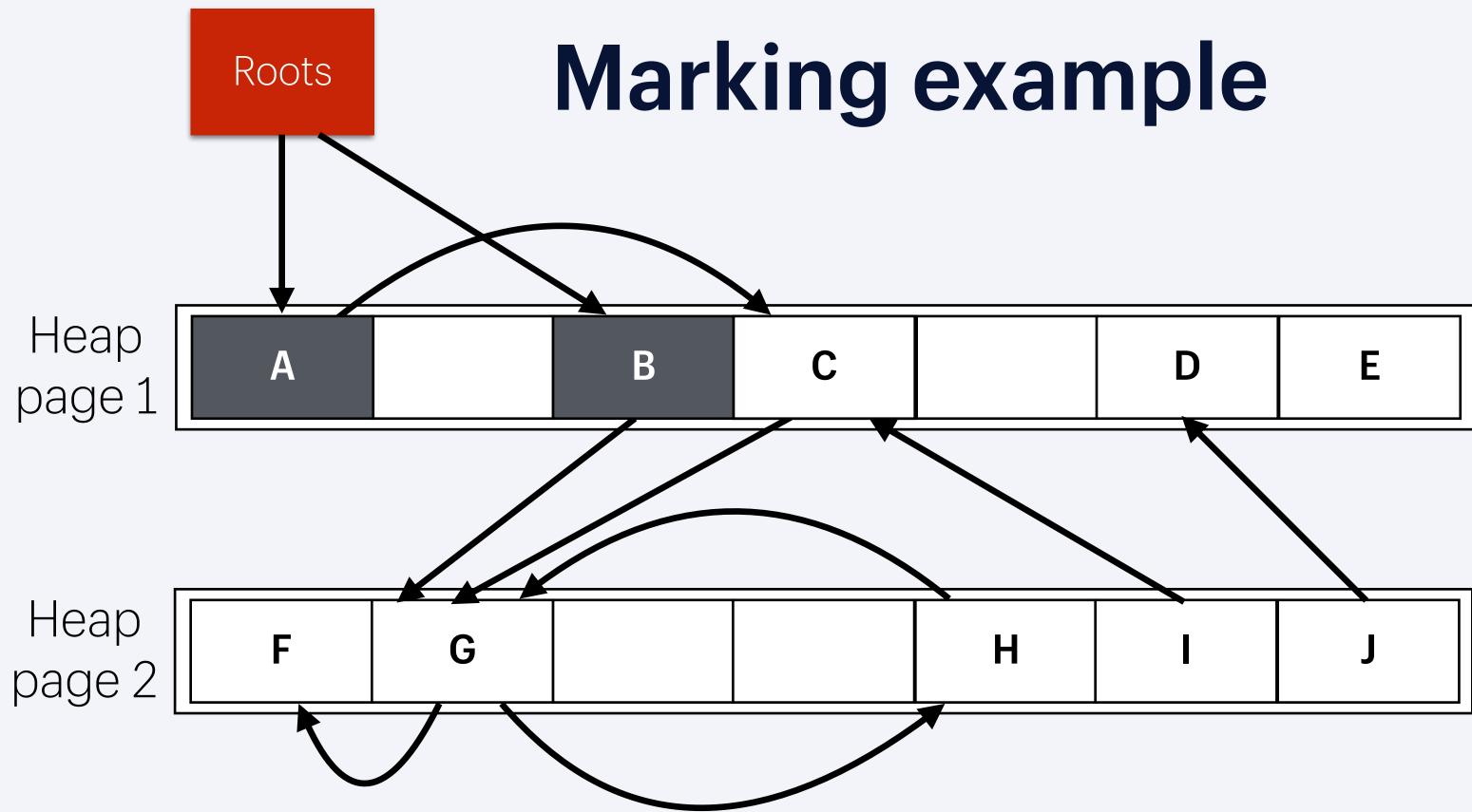
### Execute Ruby code

Mark

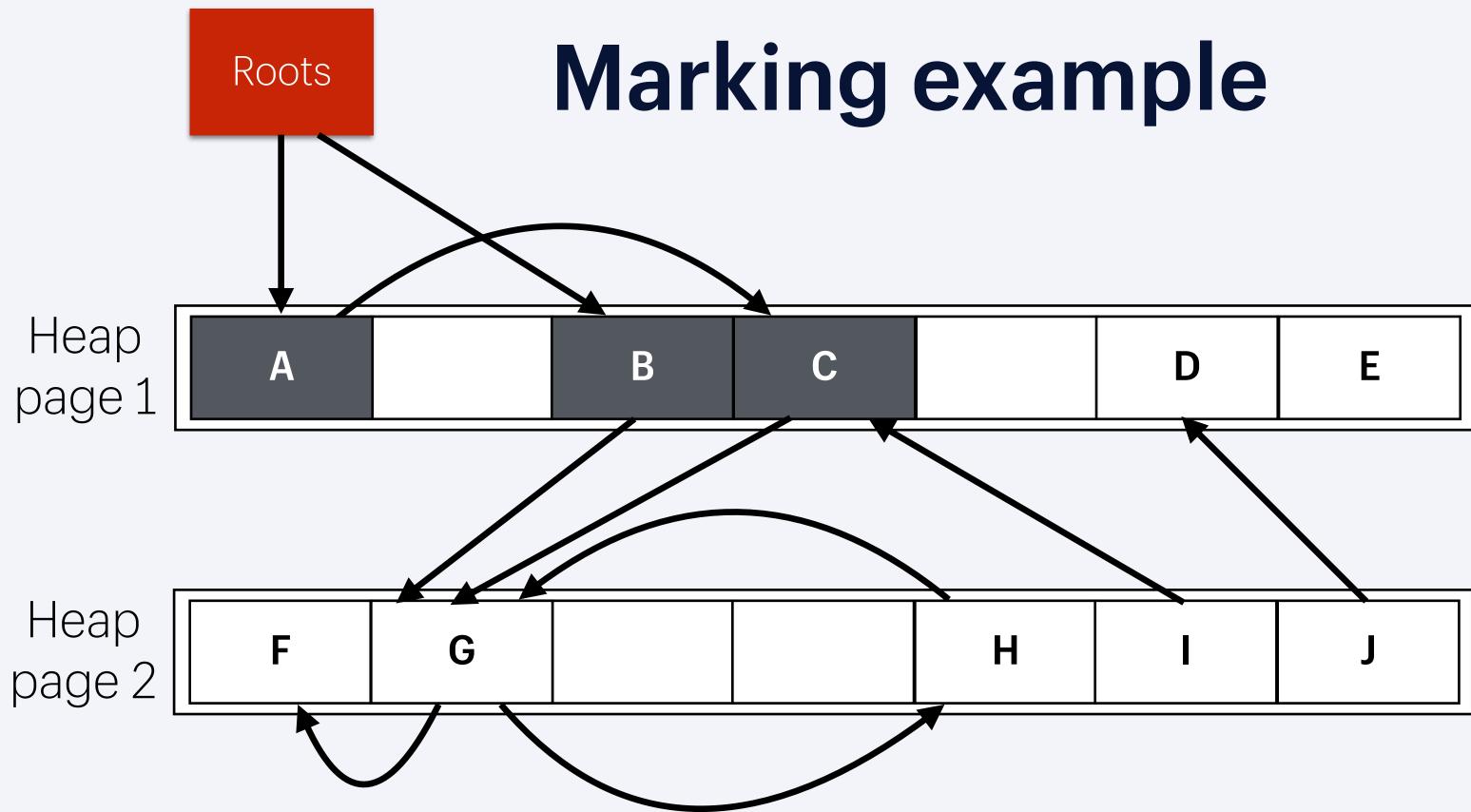




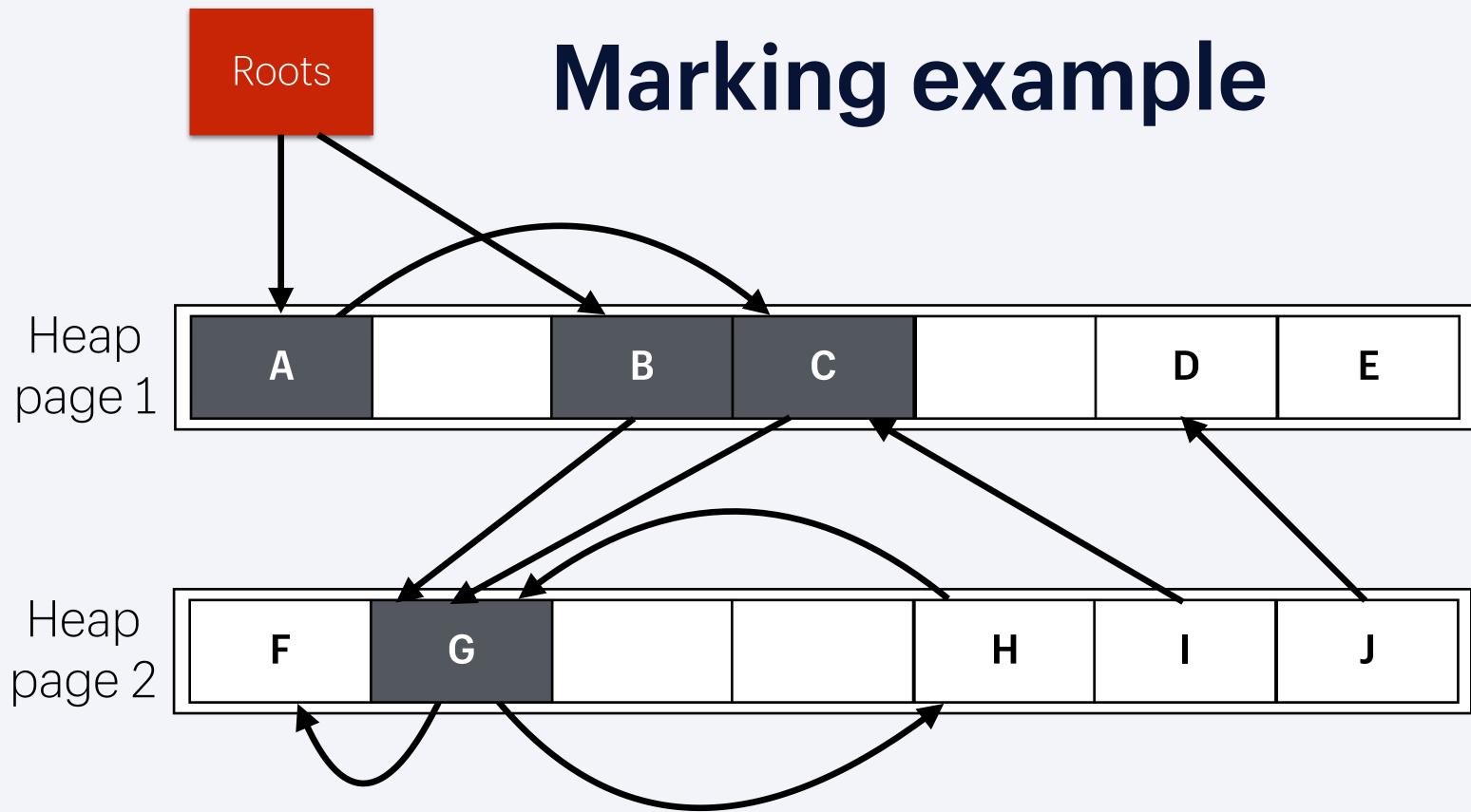




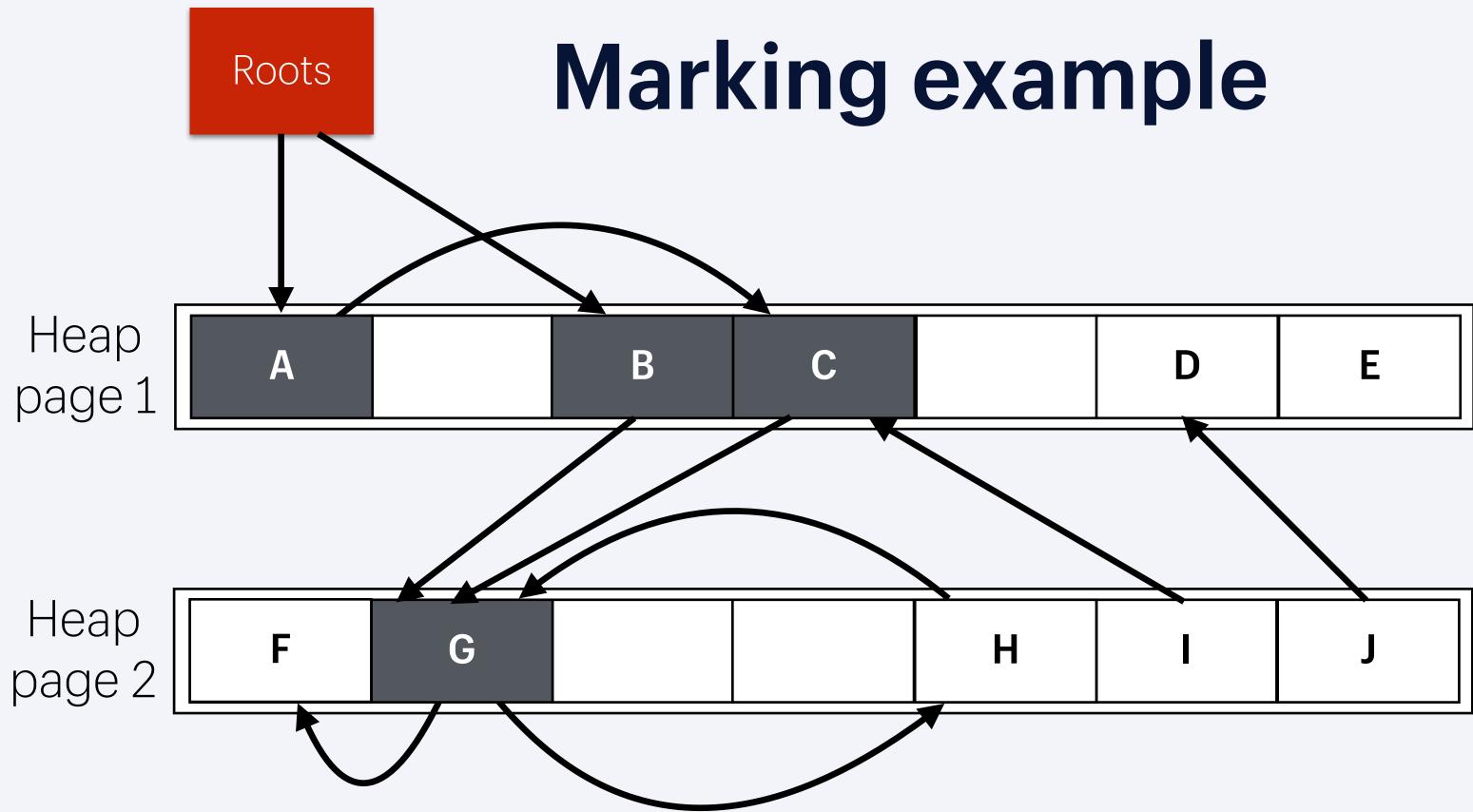


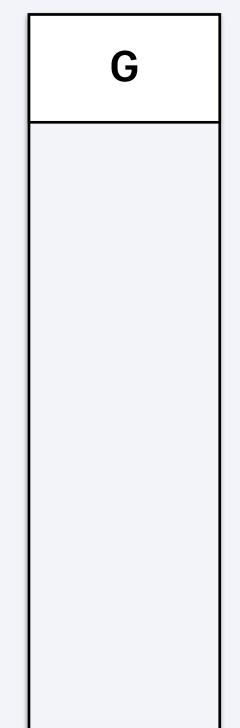


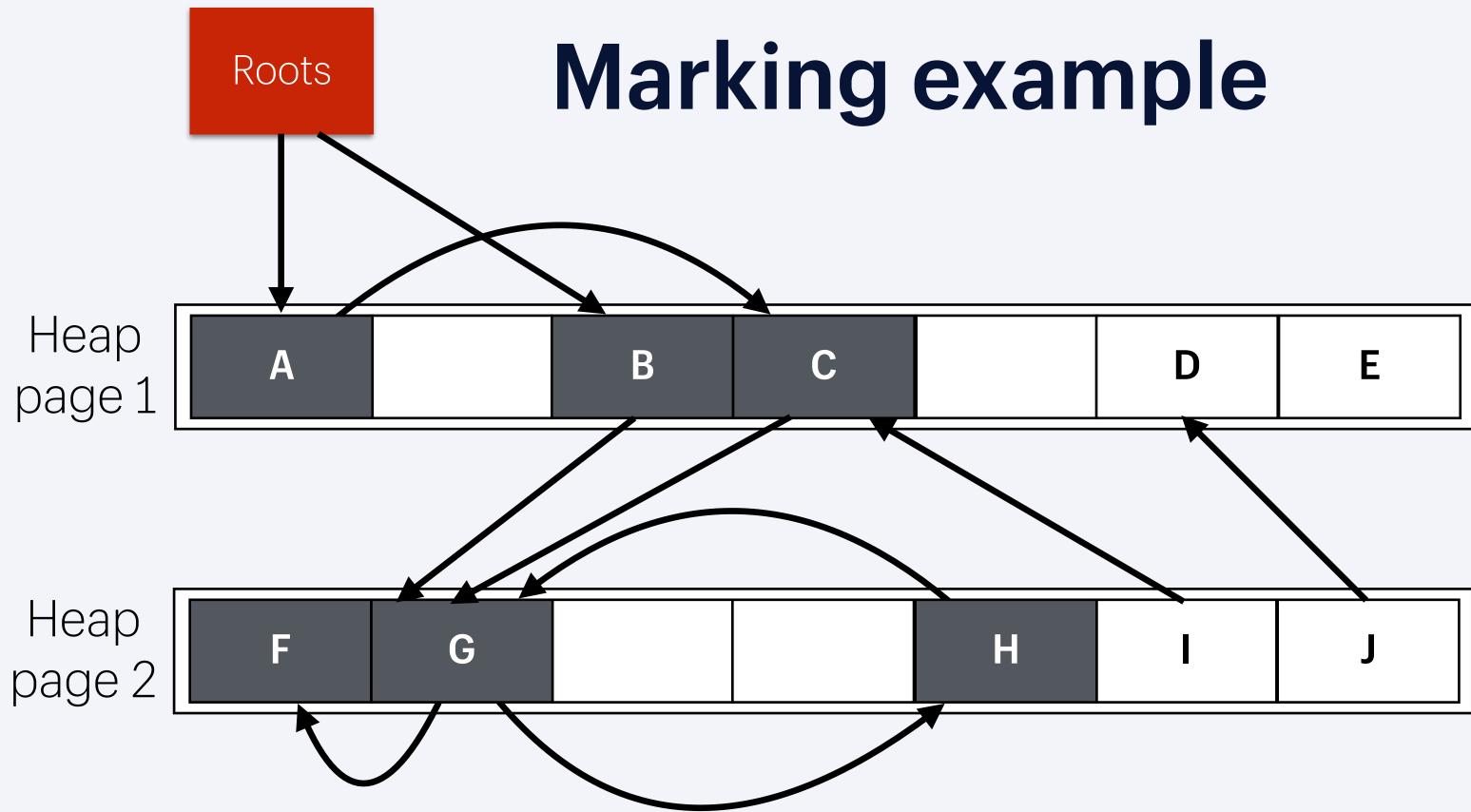




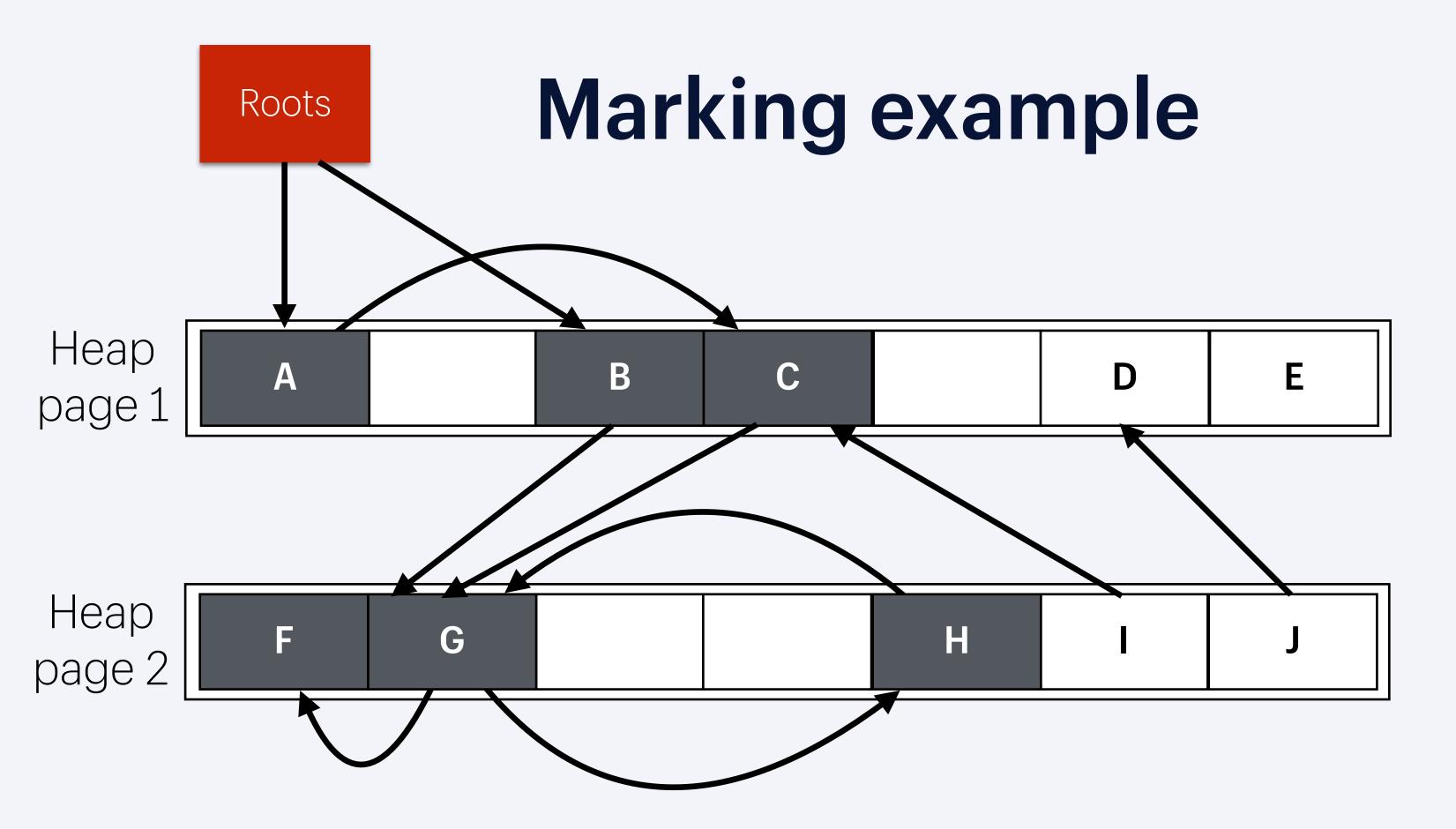








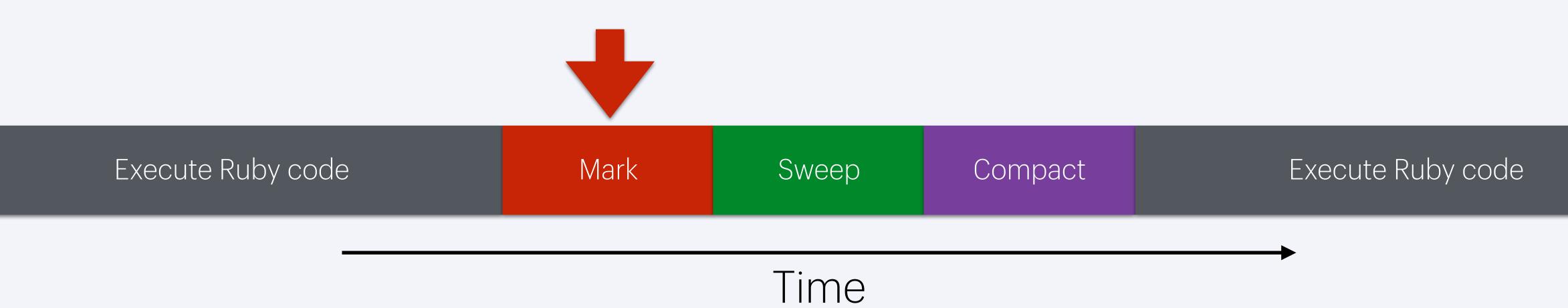


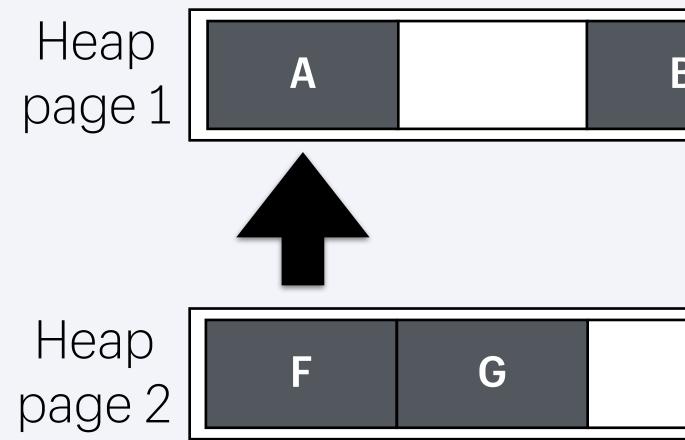




# Sweeping phase

- Marked objects = live objects
- Unmarked objects = dead objects
- Scan pages and free objects that are not marked





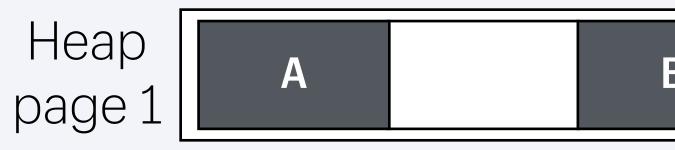
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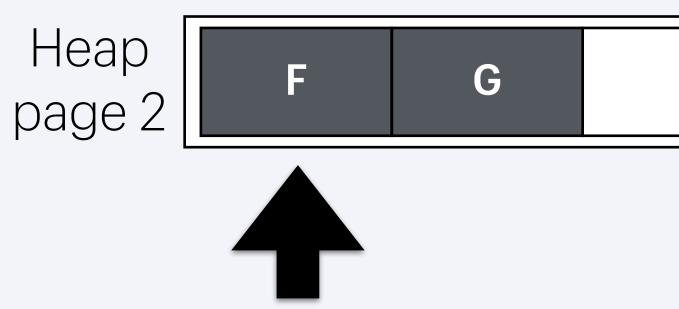
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B	С		Ε
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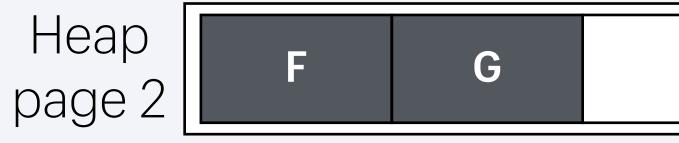




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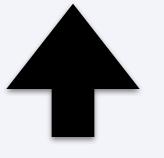
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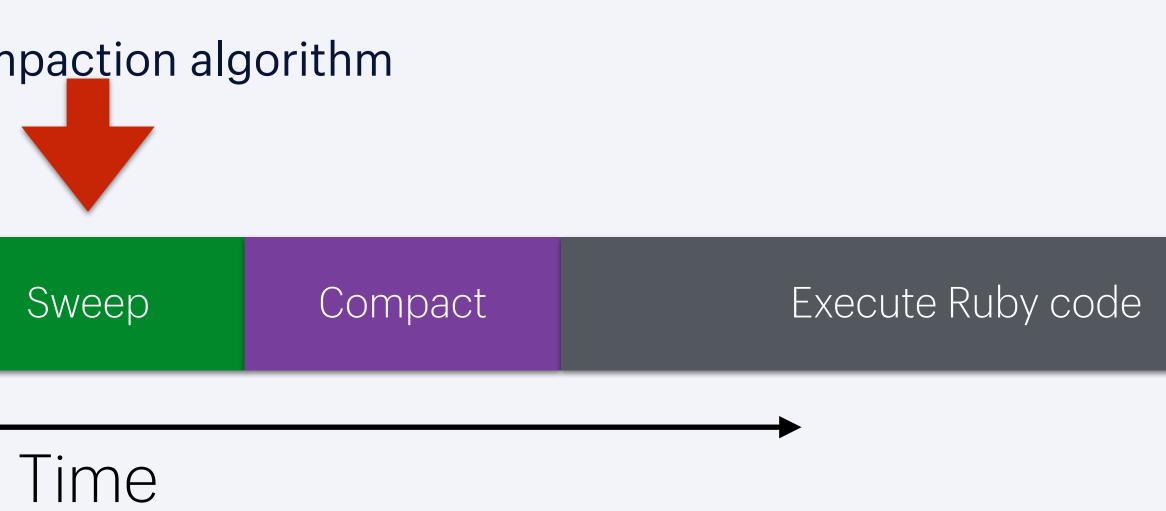
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## **Compact phase**

- Optional phase & turned off by default
- Move objects to compact the heap
- Can reduce memory usage
- Ruby uses a Two-Finger compaction algorithm

Execute Ruby code	Э	Mark	
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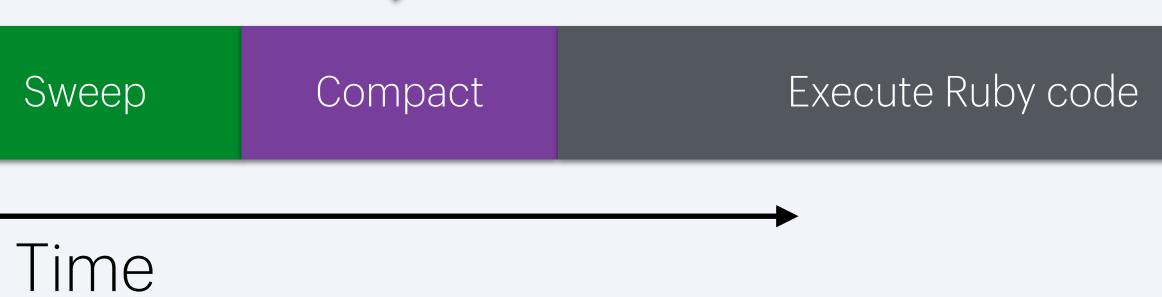


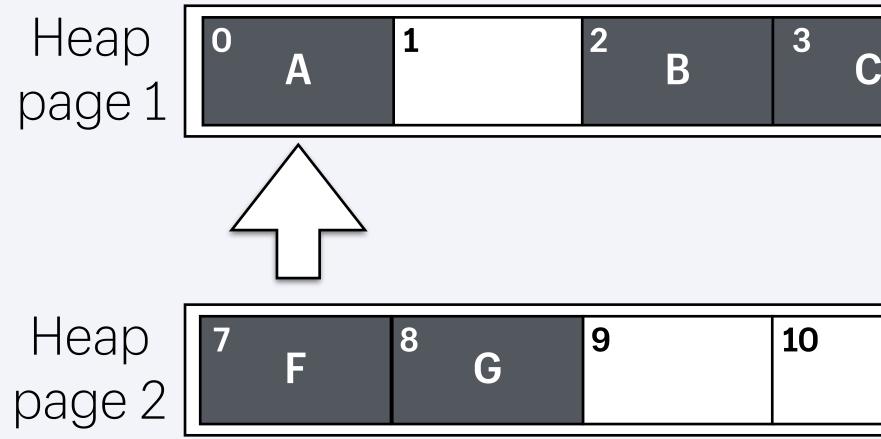
# **Compaction algorithm**

- Compact step:
  - Two cursors: compact and free
  - Free cursor moves forward and compact cursor moves backward
- Update reference step: update pointers for all objects

Execute Ruby code	2	Mark	
			-

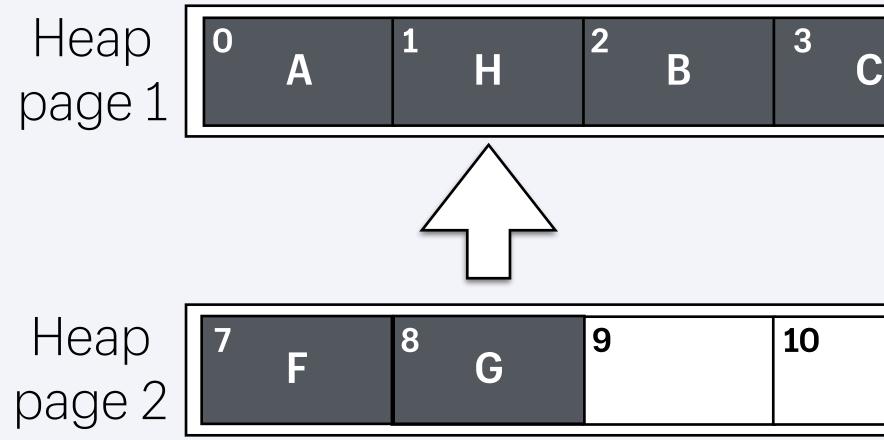






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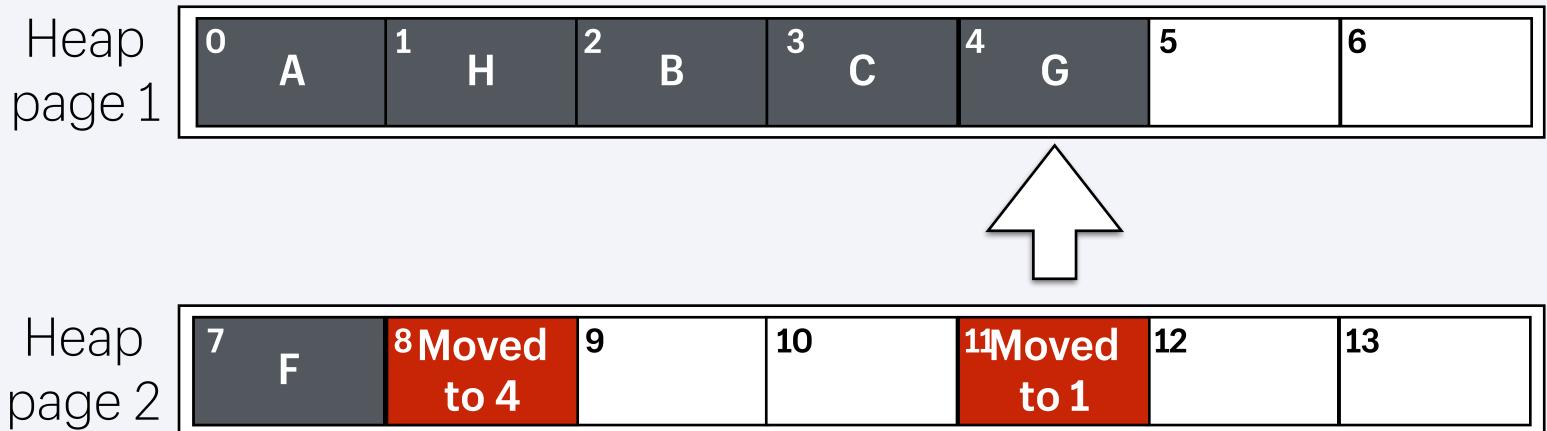
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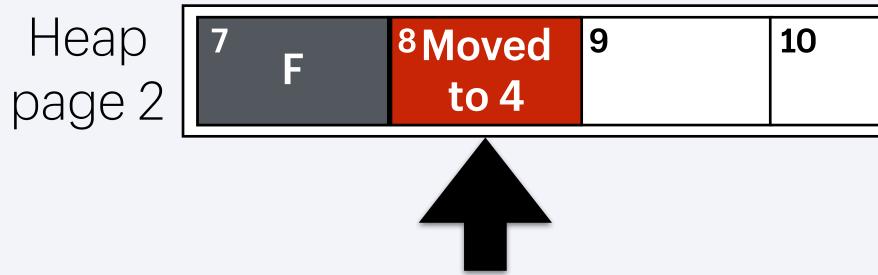


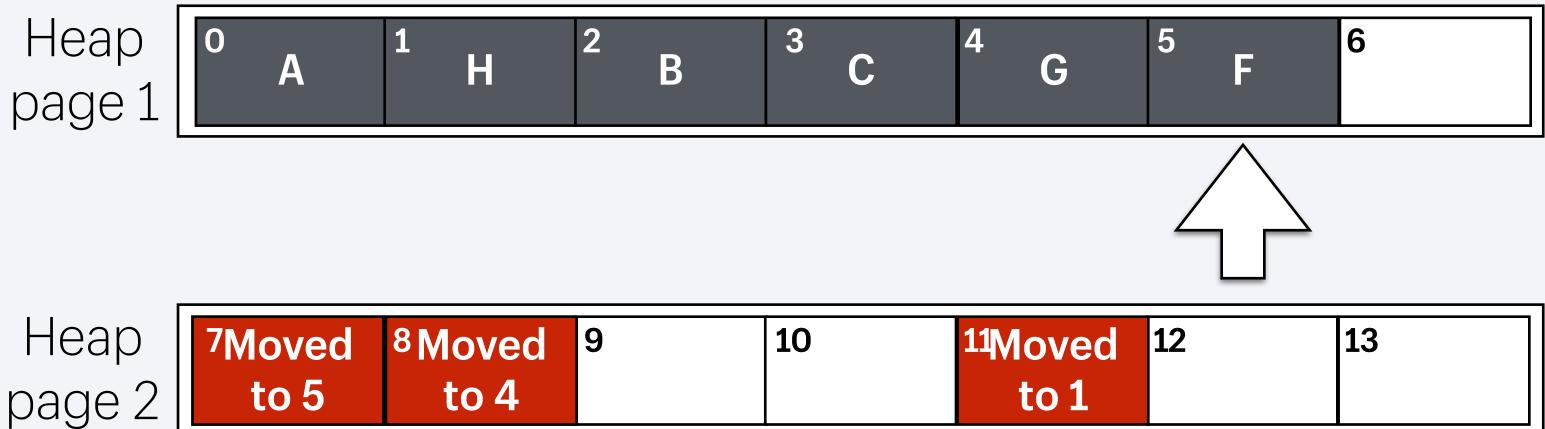
4	5	6	
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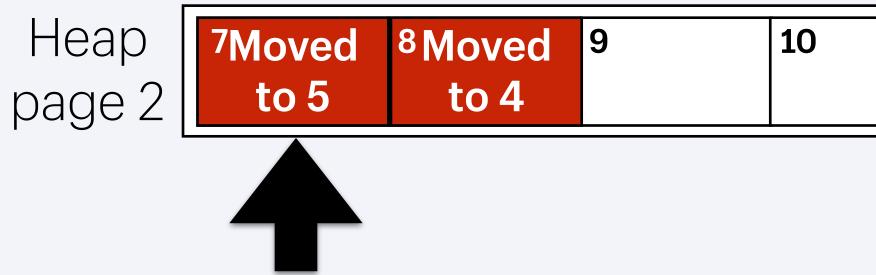
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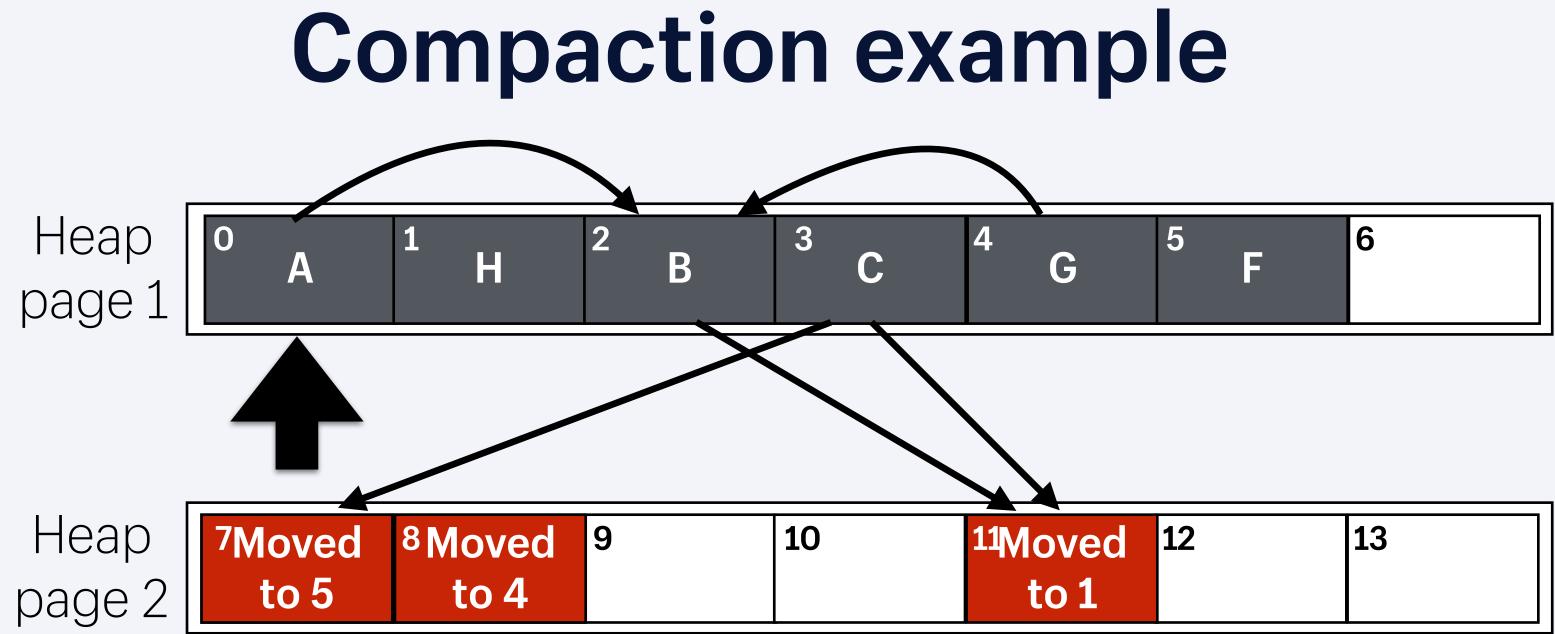


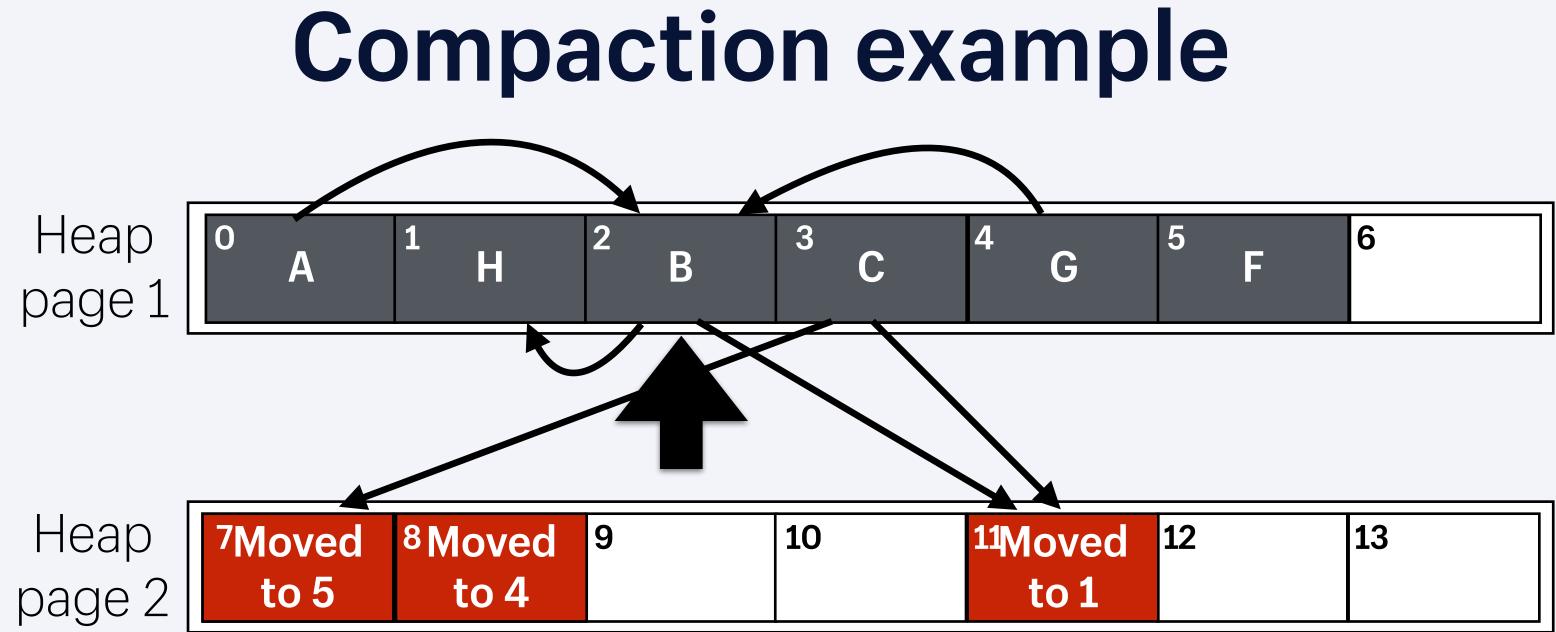


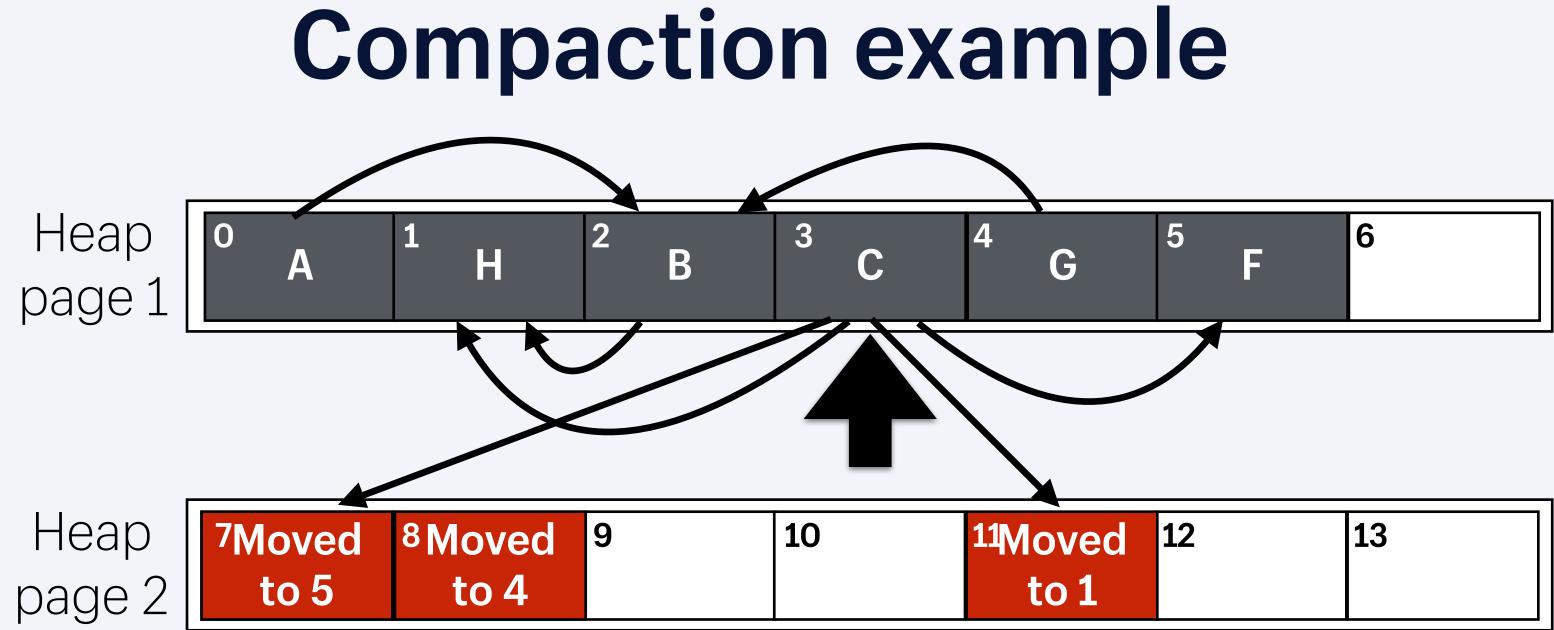


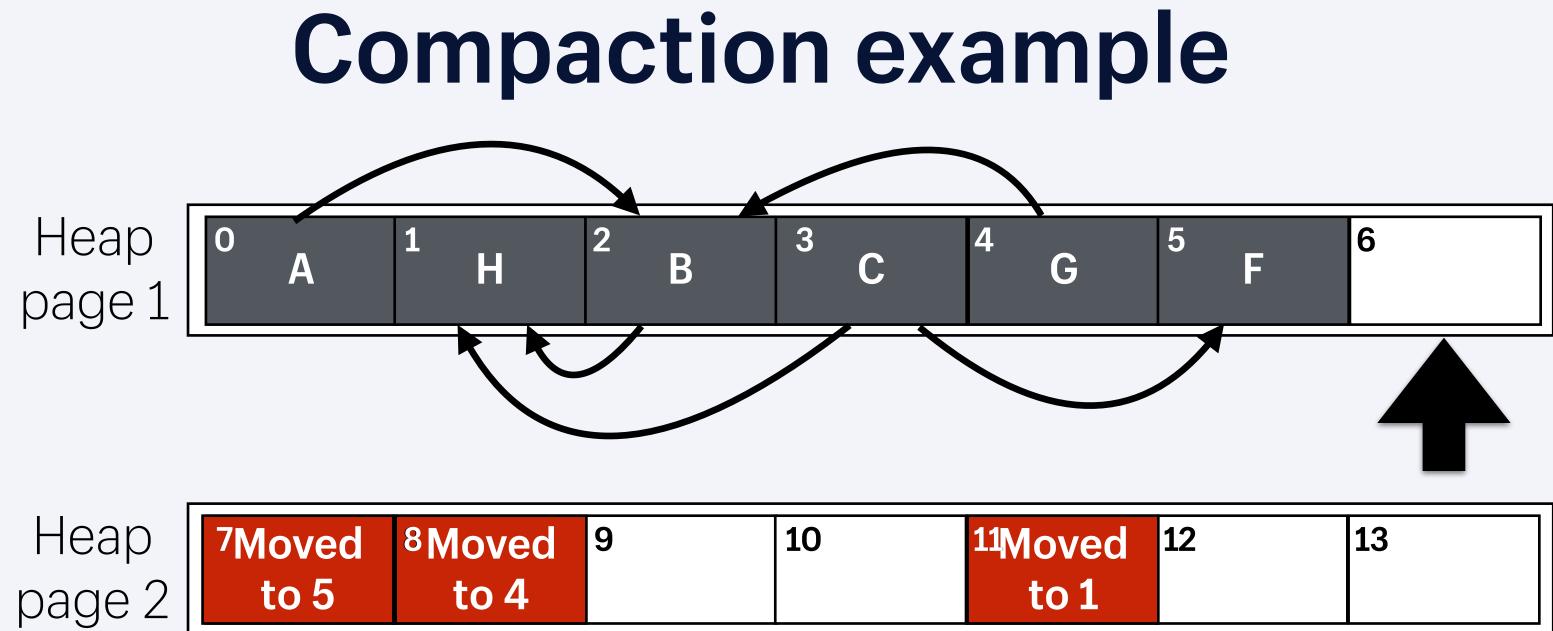












# Large objects on the heap

## Two different categories of Strings

< 24 bytes</li>

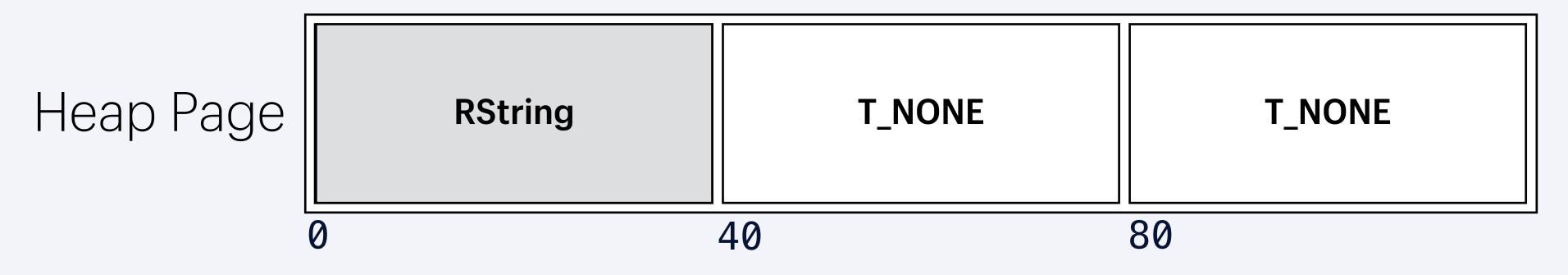
"Hello, World"

- > 24 bytes
  - "Hello RubyKaigi, thanks for having us"

12 bytes

37 bytes

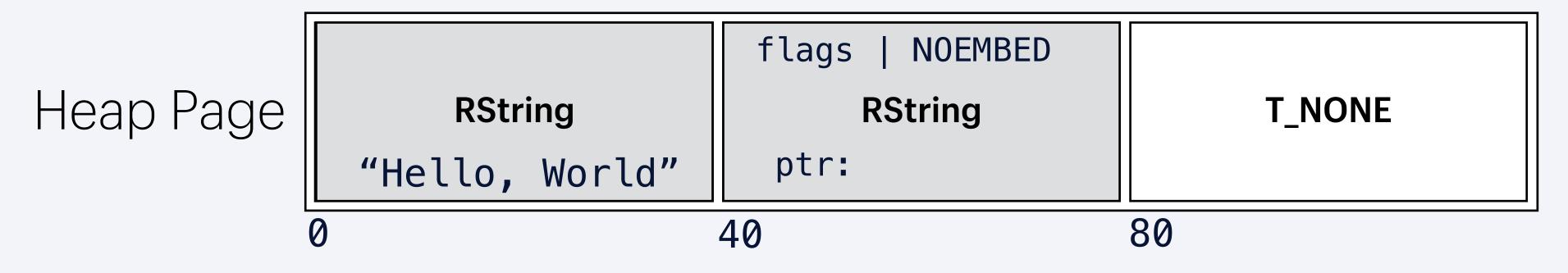
## Allocating an embedded string

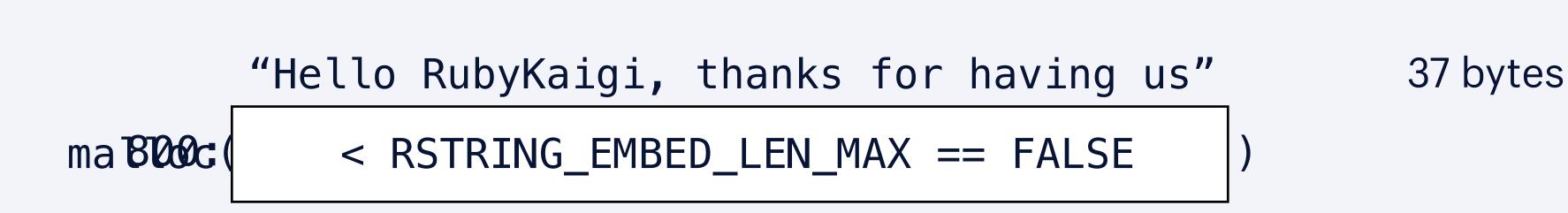


"Hello, World" < RSTRING\_EMBED\_LEN\_MAX == TRUE 12 bytes

"Hello RubyKaigi, thanks for having us" 37 bytes

### Allocating a heap allocated string





### Summary: What we've learned

- How Ruby lays out it's memory.
- How large and small objects are allocated
- What garbage collection is for, and how it works

# What problems are we trying to solve?

### Bottlenecks in the heap

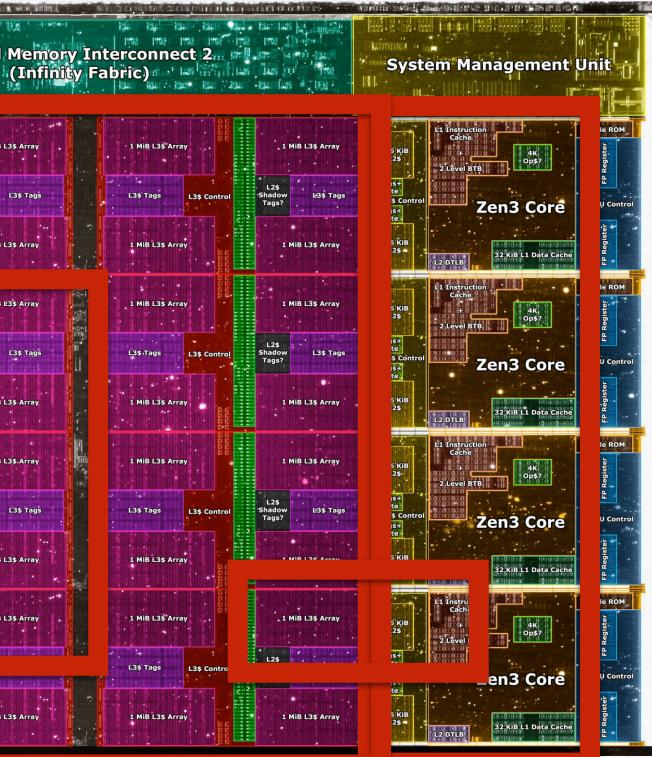
- Pointer indirection causing poor cache locality
- Performance and memory overhead caused by malloc

### **CPU caches**

- Memory in the system lives in many levels:
  - Level 1 cache: on CPU core, very fast
  - Level 2 cache: beside CPU core, slightly slower
  - Level 3 cache: shared between CPU cores, slower
  - Main memory: off CPU, very slow

### **CPU caches**

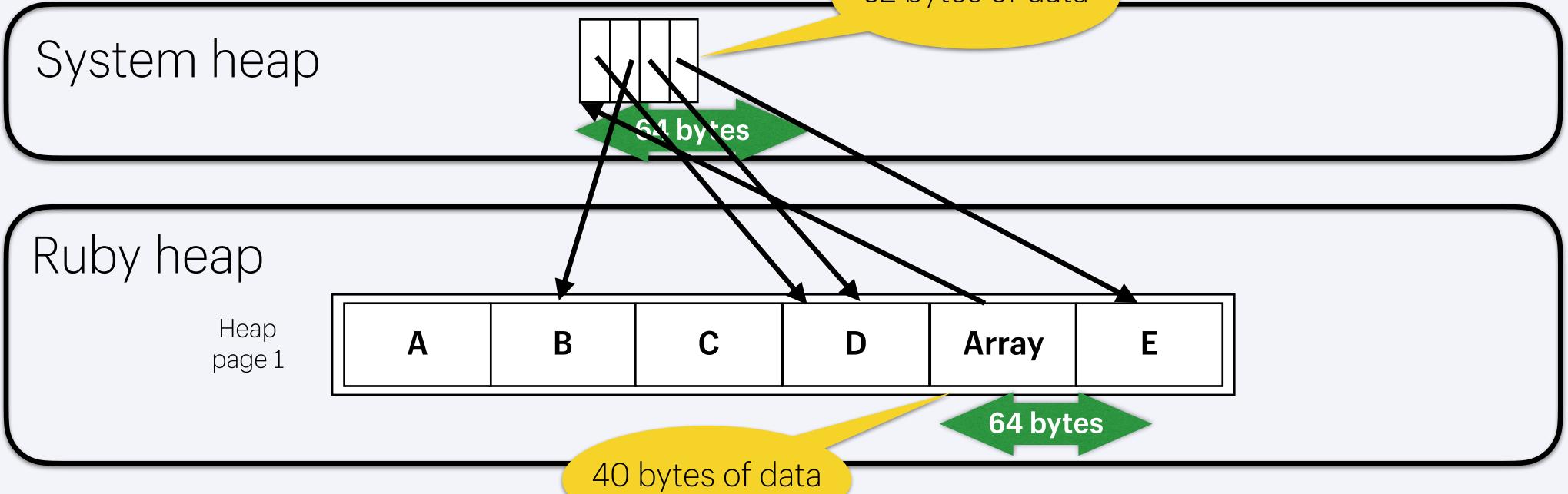
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### **CPU cache properties**

- CPU caches stores data fetched from main memory
  - CPU caches store a cache line at a time (64 B on x86)
  - Old cache entries are evicted to make space for new entries
- Cache hit: data exists in cache, no fetch from main memory required
- Cache miss: data does not exist in cache, need to fetch from main memory

### Ruby cache performance





32 bytes of data

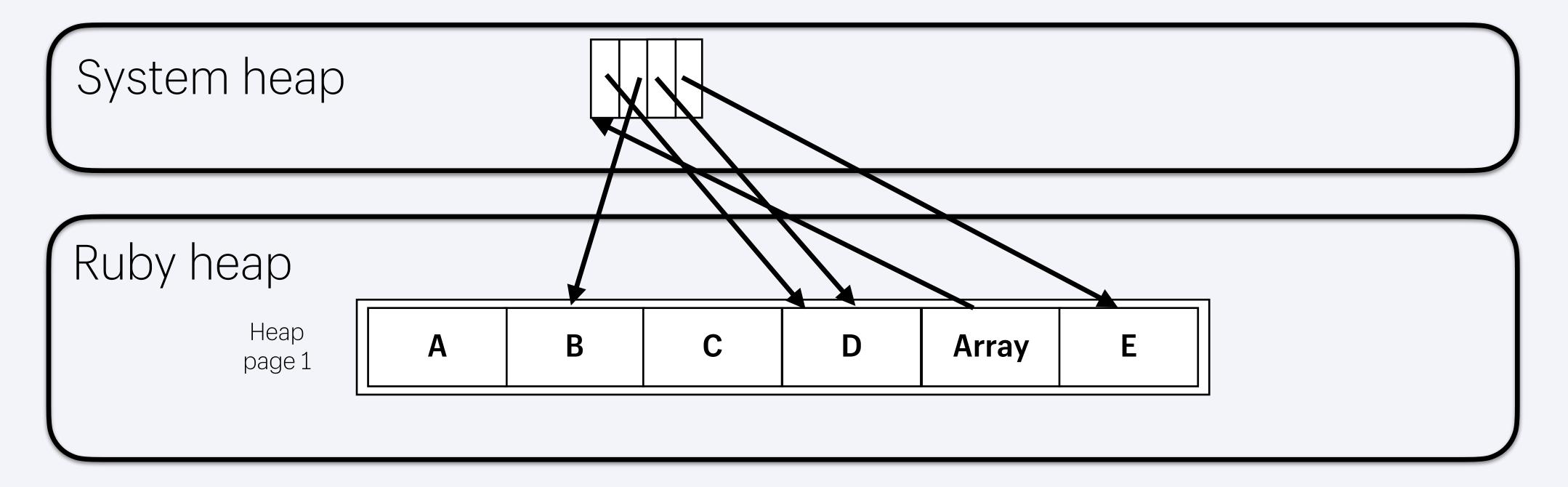
### **Overhead of malloc**

- malloc has performance overhead
- malloc stores additional metadata, increasing memory consumption
- Ruby 2.6+ introduced a second heap called the "transient heap" used to reduce the number of malloc calls
  - Increased performance in some benchmarks by 50%

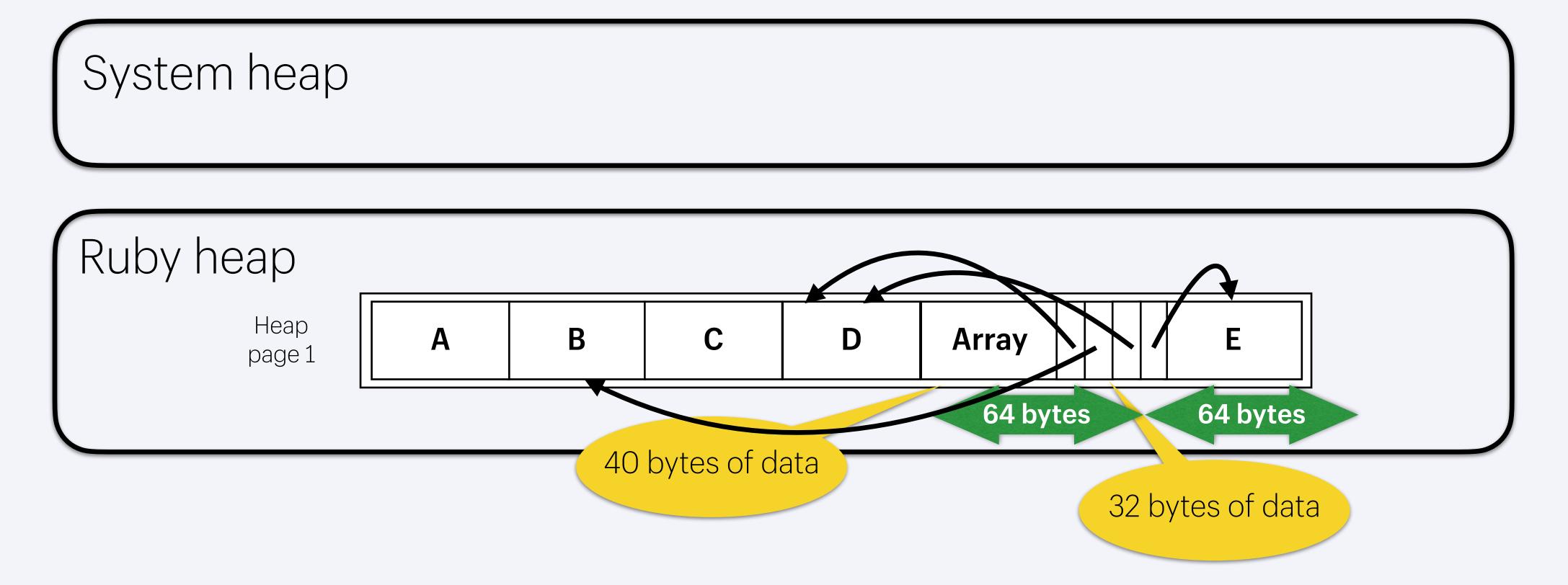
### The Variable Width Allocation project

- Extend Ruby's garbage collector to allow dynamic sized allocation
- Data will be allocated following the object RVALUE
- Reduce the number of malloc calls

### Variable Width cache performance



### Variable Width cache performance



# Where are we today?

ruby/ruby

#### #4391 **Move C heap** allocations for **RVALUE** object data into GC heap



Q 0 comments ♀ 1 review ± 7 files +452 -57 ■ ■ ■ ■ eightbitraptor • April 20, 2021 -O- 3 commits

Move C heap allocations for RVALUE object data into GC heap by eightbitraptor · Pull Request #4391 · ruby/ruby

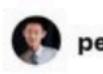
GITHUB.COM



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ruby/ruby

### #4680 Variable Width **Allocation Phase II**



peterzhu2118 • July 26, 2021 -0- 2 commits

#### Variable Width Allocation Phase II by peterzhu2118 · Pull R...

Ticket: https://bugs.ruby-lang.org/issues/18045 Feature description Since merging the initial implementation in #1757...

github.com





### \* on production workloads

export cflags
./configure
make
make
install

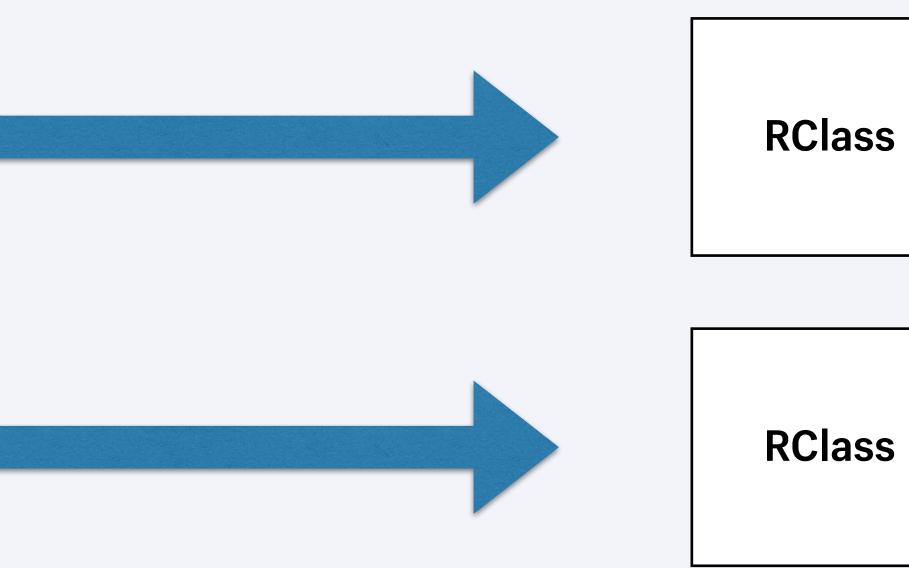
### export cflags="-DUSE\_RVARGC=1"

### **RClass allocation**

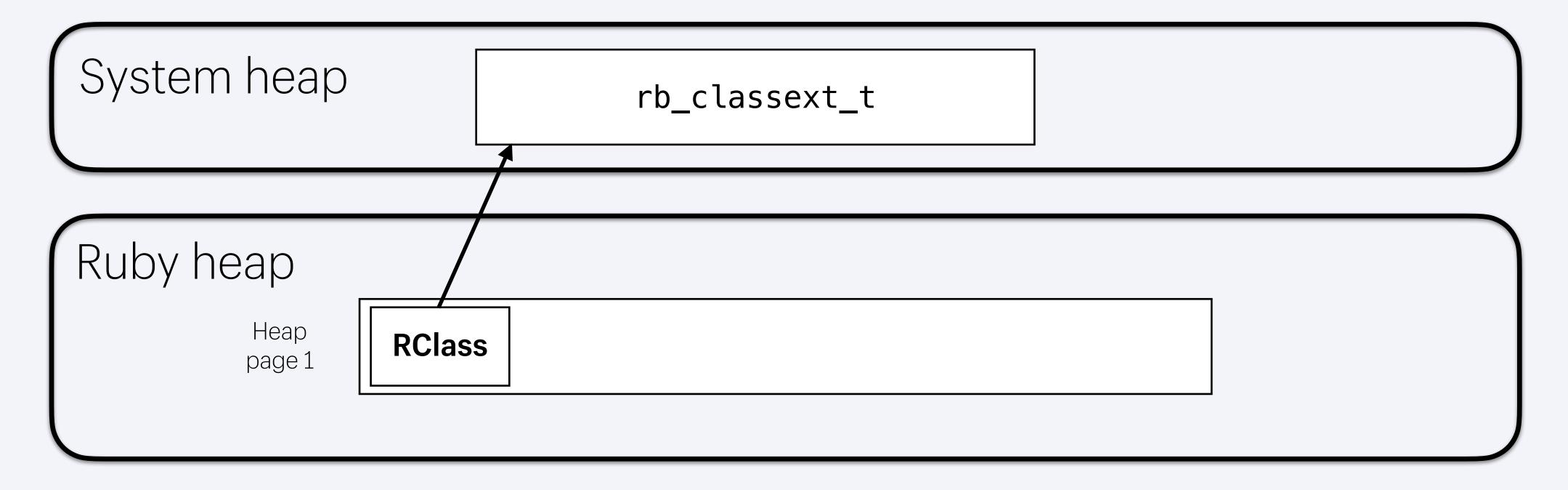
#### class MyNewClass end

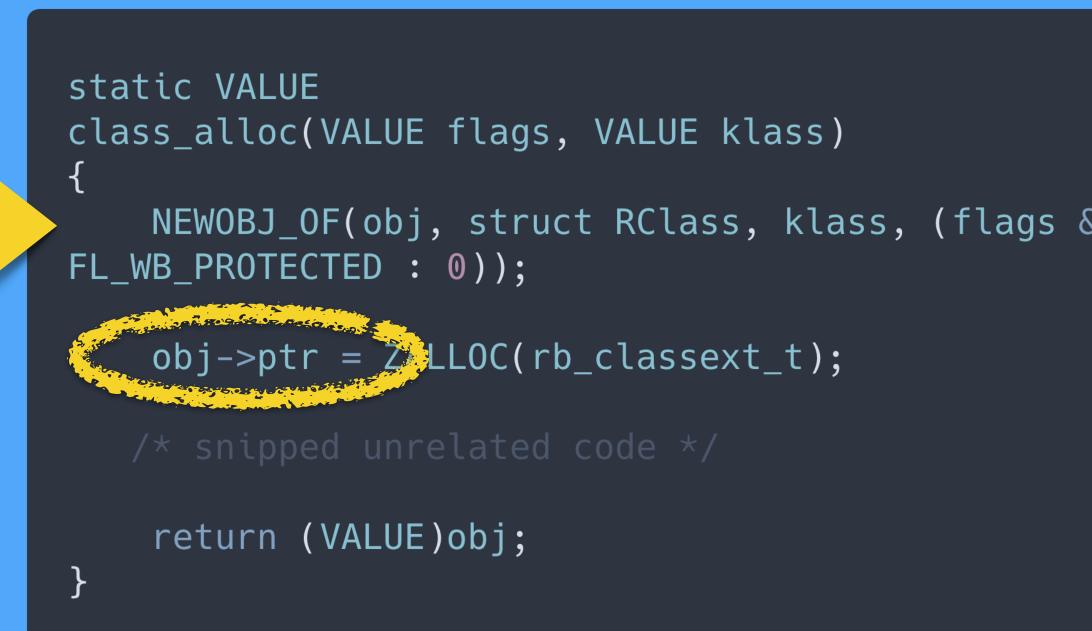
Class.new(Object)





### **RClass Allocation**





#### NEWOBJ\_OF(obj, struct RClass, klass, (flags & T\_MASK) | FL\_PROMOTED1 | (RGENGC\_WB\_PROTECTED\_CLASS ?



#### static VALUE class\_alloc(VALUE flags, VALUE klass) {

payload\_size = sizeof(rb\_classext\_t);

RVARGC\_NEWOBJ\_OF(obj, struct RClass, klass, (flags & T\_MASK) | FL\_PROMOTED1 | (RGENGC\_WB\_PROTECTED\_CLASS ? FL\_WB\_PROTECTED : 0), payload\_size);

obj->ptr = (rb\_classext\_t \*)rb\_rvargc\_payload\_data\_ptr((VALUE)obj + rb\_slot\_size());

return (VALUE)obj;

}



### Variable Width Allocation

- RVARGC\_NEWOBJ\_OF called with a desired payload size
- Object is allocated in the appropriate size pool
- Pages in size pools have different slot sizes
- Slots of size pools have powers of 2 multiples of RVALUE size

### Size pools

Size pool 0 (Slot size: 40B) Size pool 1 (Slot size: 80B) Size pool 2 (Slot size: 160B) Size pool 3 (Slot size: 320B)

### Size pools

- Allocating a class requires 40B + 104B = 144B
  - 144B = 3.6 x RVALUE

Size pool O (Slot size: 40B)

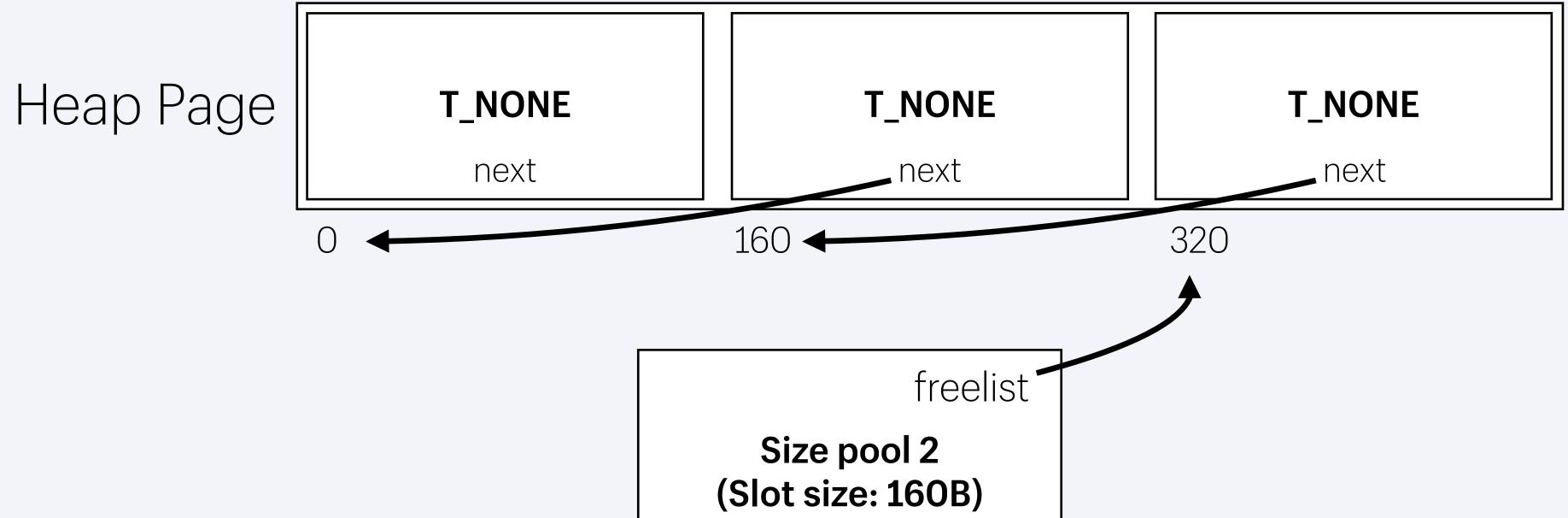
Size pool 1 (Slot size: 80B)



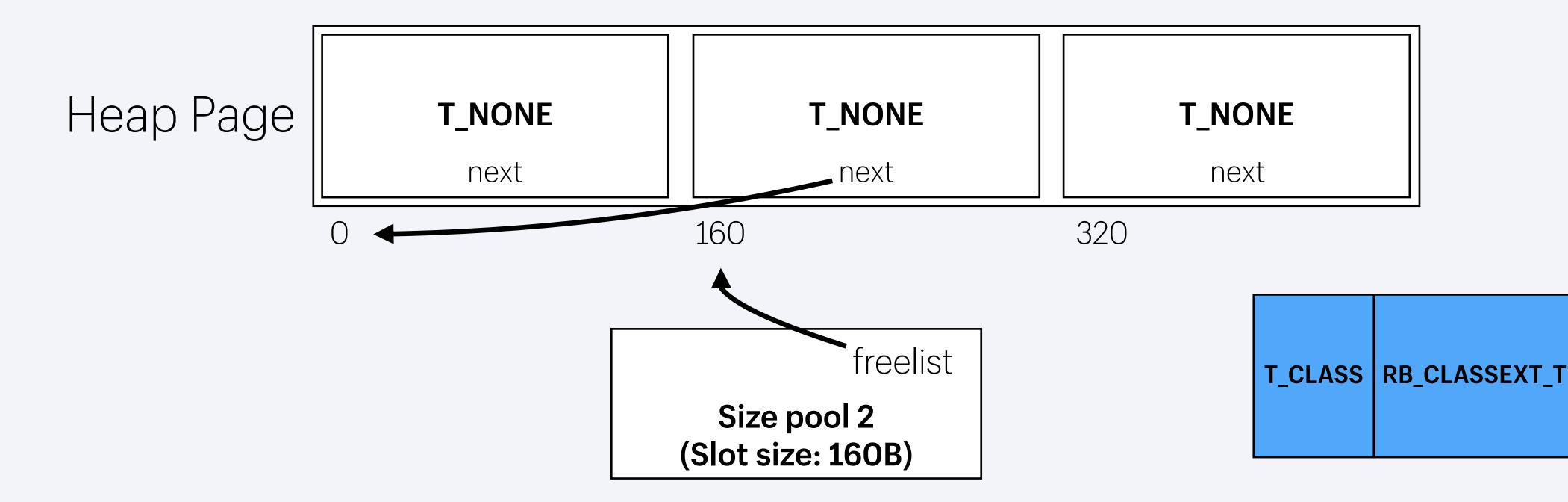


Size pool 3 (Slot size: 320B)

### Allocation



### Allocation



# Benchmarks



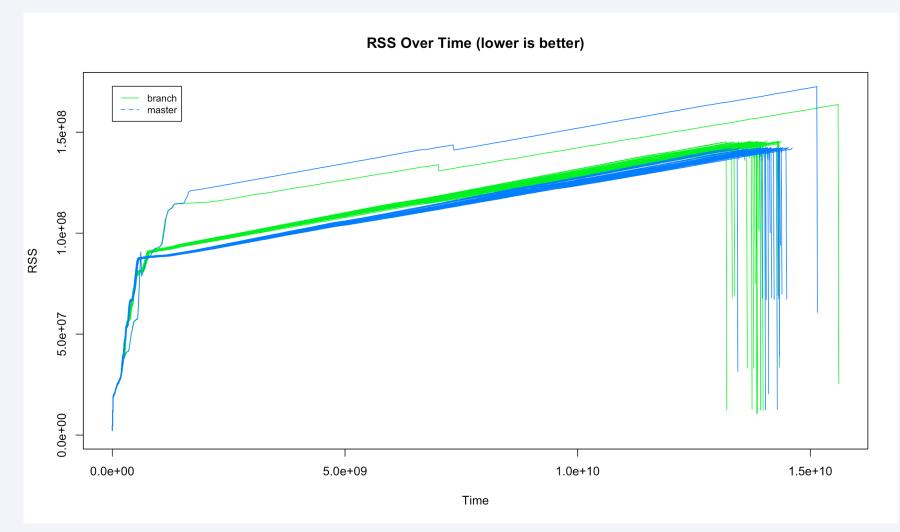
### Methodology

- Benchmarked on bare-metal Ubuntu machine on AWS
- and jemalloc allocators
- See ticket for more detailed results and analysis: https://bugs.ruby-lang.org/issues/18045

• railsbench and rdoc generation was benchmarked using the glibc

### railsbench

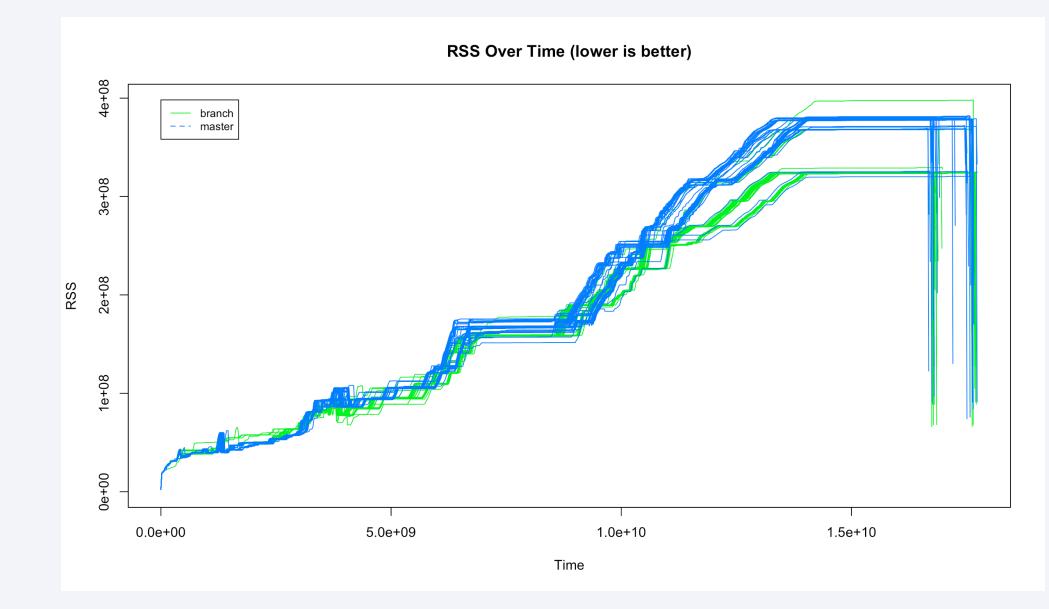
- No significant performance change when using glibc
- 2.7% faster when using jemalloc



## • 2% higher max memory usage when using glibc and jemalloc

### rdoc generation

13% lower memory usage th jemalloc



#### • 13% lower memory usage than master when using glibc and

### Liquid & optcarrot benchmarks

No significant performance difference beyond margin of error

# Limitations and future plans

### VWA everywhere

- Currently only classes are using Variable Width Allocation
- Add support for arrays and strings

sing Variable Width Allocation strings

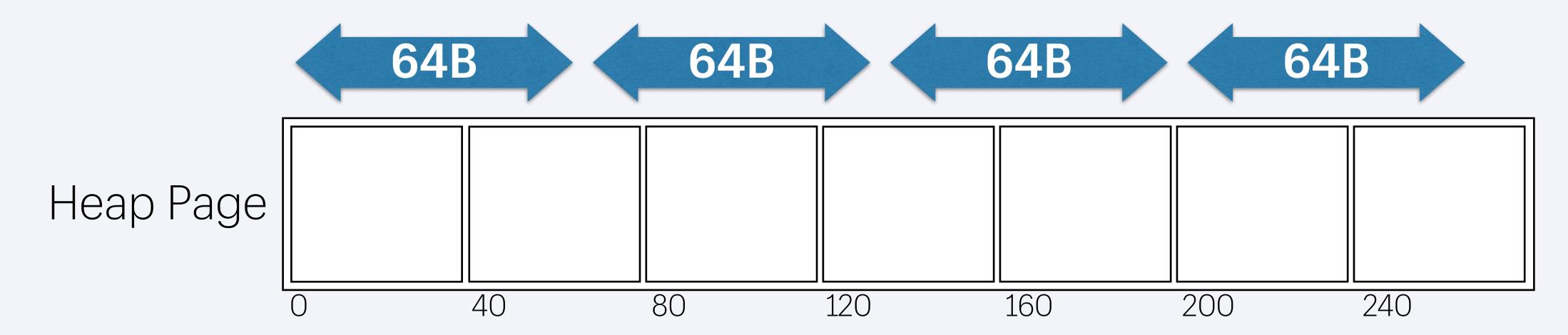
### **Resizing objects**

- Arrays and strings can resize upwards
- Difficult problem to tackle
- advantage of compaction to move resized object

• One idea: allocate extra space in a larger size pool and take

### Shrinking RVALUE

- We'd like to shrink RVALUE from 40B to 32B
- Align on 64B cache line boundaries



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- Align on 64B cache line boundaries

